Reader Collision Avoidance Scheme for Mobile RFID–Sensor Integrated Networks

Doo-Hyun Ko*, Song-Min Kim*, Sang-Bin Lee, Sun-Shin An★

Abstract

In recent years, one of alternatives for constructing RFID networks that provide mobile services is using wireless sensor networks (WSN) to enhance network capacity, utility and scalability. Due to absence of compatible reader anti-collision control and channel capture phenomenon, the medium access control protocols as used in the RFID networks lead to reader collision and starvation problem. In this paper, we develop a MAC protocol which is called Enhanced Collision Avoidance MAC (ECO) to avoid reader to reader collisions in an integrated RFID network. ECO is a CSMA-based MAC protocol, and operates on integrated nodes which consist of a RFID reader and a mote. Performance evaluation shows superior results to pure-CSMA protocols under dense deployment environments, both in number of failures and in throughput.

Key words: MAC, WSN, RFID, Network

I. Introduction

As electronics and communication technologies advance, wireless communication system has been accomplished incredible growth, and major devices also are became smaller and smarter. This advance will elevate computing systems, and it will be not easy to find the ubiquitous system in daily lives. One of the key technologies for ubiquitous computing is radio frequency identification (RFID) which is a data-collection technology that uses radio waves to remotely store and to retrieve data to identify various objects. It provides companies with up-to-the-minute supply chains visibility, including inventory, logistics and freshness dates. Recently, RFID is increasingly being used in many applications and services such as inventory management, object tracking, commuting check etc. For example, KLM/Air France implemented RFID-enabled baggage tags inside Amsterdam Airport, and Google India a division of search engine company Google, is using an RFID system to track up to 100,000 of its assets at its offices. However, as RFID systems and mobile RFID services increases, data management is more important and complicated. Consequently, RFID systems should be managed with networks for efficient data collection and transmission. One of alternatives for constructing RFID networks is using wireless sensor networks (WSN). Although monitoring of the surrounding physical environment is indispensable in the management of environment-sensitive objects, existing RFID networks and services lack this capacity. The environmental information of RFID networks, such as temperature and humidity is also not available. In this cases, the surrounding conditions can be

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learned through wireless sensor networks (WSNs), which are wireless networks consisting of spatially distributed autonomous devices using sensors to cooperatively monitor physical or environmental conditions, such as temperature, sound, vibration, pressure, motion or pollutants, at different locations. WSNs were originally developed for military applications such as battlefield surveillance. However, WSNs are now used in many commercial areas, including environment and habitat monitoring, healthcare applications, home automation, and traffic control. Capacities of WSNs are also very powerful as a network, because WSNs not only support data collection but also provide routing, aggregation and analysis of information. Furthermore, WSNs are considerably attractive when constructing RFID networks due to their high compatibility with other networks. To complement defects of RFID networks, they will ultimately need to combine with WSNs. In SARIF [3], it illustrates well a framework for integrating wireless sensor and RFID networks. They proposed a novel framework, SARIF (sensor and RFID integration framework), that integrates RFID networks and WSNs for environment-sensitive object tracking and management, and proved that SARIF can achieve energy efficiency by load balancing. But integrated RFID still remain to solve several problems. One of them is the collision problem that is arisen between tags and readers. Currently, common solutions exist for tag-tag collision, and reader-tag collision has been also studied in the EPC-Global Class1 Gen1 and Gen2 standards for UHF readers. However, although there are some algorithm for reader-reader collision such as Colorwave [4], PULSE [5] and IRCD [6] etc, these algorithms are not solutions for sensor and RFID integration framework. Consequently, we focus on reader-to-reader collision avoidance scheme for integrated RFID-sensor networks.

In this paper, we propose an integrated RFID-sensor networks architecture that enables an application to be designed flexibly while hiding the details of the RFID networks and WSNs. And we discuss a scheme to avoid collisions in multi reader scenarios.

II. System Design for Mobile RFID-sensor Networks

In this section, we propose an architecture of integrated RFID-sensor networks in a mobile environment. Our network structure can provide a remote monitoring service and a mobile terminal service. Currently, it is possible to get information of products from mobile terminals. This service is using cellular networks, in which a reader is embedded in a mobile phone, and if the reader read a tag, the mobile can get information from tag’s data via cellular network. However, it is the restrictive service which is for only a mobile user. The RFID system should be used various objects and provided with useful information. For example, if one consumer reads attached tags to buy fresh foods, it should be provide not only cost and date of manufacture but also temperature data in the storage. RFID system should also provide sensing information which is requested to control resource management. Consequently, by integrating RFID systems with WSNs, we can build an object tracking system and management system that can provide richer information about the environments of objects.
is transmitted to the sink via sensor networks, where the sink node collects not only sensing data but tags. The collected data is sent to server or application through core network. When clients request information to server, server provides appropriate information. This network structure can provide not only real-time service but also tracking and mobile services.

II- Hardware Configuration

The integrated RFID-sensor module is consisted of three parts which are RFID reader (900MHz), micro-controller and RF transceiver (2.4GHz). We use the SkyeModule M7 (RFID reader module) which is globally compliant UHF module for embedded RFID. The read range of this reader is up to 1 meters with external antenna. It can communicate with a host micro-controller over the UART, SPI and I2C interfaces. The interface between microcontroller and reader parts is UART bus. The reader module is capable of responding to ASCII and binary commands sent by the host micro-controller. It can select, read and write tags. The host controller can also read and write the reader’s memory and system registers to put the reader in low power sleep mode and to wake it up from sleep. The small footprint and low power requirement makes it suitable to integrate RFID and sensor networks. A used RF transceiver is CC2420 which is known well in sensor network area. The CC2420 is a true single-chip 2.4 GHz IEEE 802.15.4 compliant RF transceiver designed for low-power and low-voltage wireless applications. CC2420 includes a digital direct sequence spread spectrum (DSSS) baseband modem providing a spreading gain of 9dB and an effective data rate of 250 kbps. It is also connected to micro-controller with SPI bus. This module is called as Ubicon. Fig 2 and 3 illustrate Ubicon and its block diagram.

In proposed network, we use the integrated module. When the number of RFID readers is increased or when the module is moving, the probability of reader collision increases. Consequently, we have designed a medium access control protocol to avoid the reader collision problem. The protocol is discussed in this section. we propose CSMA-based reader control scheme for communication with WSNs, and we call this scheme as enhanced collision avoidance MAC (ECO).

III. Enhanced Reader Collision Avoidance Scheme for Mobile RFID-sensor Networks

Simultaneous transmission in RFID systems lead to collisions as the readers and tags typically operate on the same channel. In this section, we will discuss about two interference types that are caused reader collisions before describing ECO.

A. Reader to Reader Interference

Reader-to-reader frequency interference, or simply frequency interference, occurs when a reader transmits a signal that interferes with the operation of another reader,
thus preventing the second reader from communicating with tags in its interrogation zone. This type of interference occurs when the signal transmitted by a reader is of sufficient strength when received at a second reader that the signal masks or jams communication from tags to the second reader. Interrogation zones need not overlap for reader-to-reader frequency interference to occur. The signal from the tag to the reader is extremely weak since tags communicate with readers by either reflecting or loading the reader’s own transmission. These signals are easily masked by transmissions from other nearby readers. For example, R₁ lies in interference region of reader R₂ in Fig. 4. The reflected signals reaching reader R₁ from tags, can easily get distorted by signals from R₂. Note that such interference is possible even when the read range of the two readers do not overlap.

B. Multiple Readers to Tag Interference

Multiple reader-to-tag interference, or simply tag interference, occurs when one tag is simultaneously located in the interrogation zones of two or more readers and more than one reader attempts to communicate with that tag at the same time. In this type of interference, each reader may believe it is the only reader communicating with the tag while the tag is in fact connected with multiple readers. In simple RFID networks, the protocol is implemented in an undesirable way that may interfere with the communicating readers’ abilities to communicate with that tag and other tags in their respective interrogation zones. In Fig. 5, the reading range of the two readers is overlap. Hence the signals from R₁ and R₂ might interfere at tag T₁. In this case, T₁ cannot decipher any query and the tag is read neither by R₁ nor by R₂. In this case, R₁ will indicate presence of two tags instead of three.

III-II Enhanced Reader Collision Avoidance MAC (ECO)

We have implemented enhanced collision avoidance MAC using the unslotted CSMA-CA algorithm. ECO is an asynchronous MAC for mobility service. We should also consider controlling both of channels when we design a
protocol, because our integrated module has two channels. Refer to fig 6, ECO arbitrates the RFID-sensor networks with two-way handshake of RTR-PTR packets. The RTR packet means that a node (Ubicon) requests to read adjacent tags. The PTR packet is response of adjacent nodes for the RTR packet. The RTR packet includes several information such as reading time, transmitter’s ID and etc. If one of Ubicons wants to read tags around, it broadcasts the RTR to other nodes. Nodes which receive the RTR packet send the PTR packet to permit reading, and then they fall in wait phase or sleep mode during reading time to guarantee reading of the RTR packet transmitter without collisions. When each node which has received the RTR broadcasts the PTR, they choose randomly one of n slots within threshold time to avoid collision of the PTR.

![Fig. 8. Readable Ubicons after RTR-PTR packet handshake](image)

Each slot is very short time compare to backoff. The node which sent the RTR, reads immediately tags, when the node receives first the PTR. If the node doesn’t receive PTR during threshold time, the node regards adjacent nodes as nothing; therefore, the node starts reading. It is illustrated in fig 7. The PTR packet is also used to give a reading chance to adjacent nodes. When the PTR is broadcasted, the nodes which received the PTR also choose randomly slot and send RTR message, if those want to read the tags. As this way, ECO increases the reading chance of nodes. Refer to fig 8, R1 that has the shortest backoff broadcasts the RTR packet after backoff time, and then adjacent R2, R3 and R4 that received the RTR packet send the PTR packets and store the ID of the transmitter in the RTR packet. Nodes (R2, R3, and R4) which received the PTR packet also store the ID from PTR, and send again RTR packets. If these nodes don’t receive any other packets or receive the PTR packet, these nodes also have simultaneously a chance which can read tags. R6 should sleep in reading time because it received the RTR from R5.

III-II-I ECO Algorithm Description

Fig 10 shows the detailed algorithm for the ECO protocol. This description illustrates that our reader (Ubicon) is how to operate in the integrated network. countRTR is zero, where the countRTR is the number of received RTR, and when the backoff time is elapsed, Ubicon broadcasts the RTR and waits until receiving the PTR or waitTerm (threshold time). And then Ubicon reads tags through the UHF channel. If Ubicon completes reading tags, it goes into sleep mode. In state of Received RTR; if already countRTR is larger than zero, Ubicon adds one in the countRTR, and waits until reading time is elapsed. Then Ubicon subtracts one from the countRTR. If the countRTR had been zero, backoff counter is stopped immediately, and broadcasts the PTR. At this time, The countRTR is set as one. Then Ubicon waits during reading time. After that, if the countRTR is zero, Ubicon goes to countRTR is zero: state to get a reading chance. If not, it goes to Received RTR: state. when the reader receives the PTR, it goes to Received PTR: state. In this state, the reader waits during backoff time, and if the countRTR is zero, goes to countRTR is zero: state to get reading chance, if not, goes to Received RTR: state.

```
Begin ECO

countRTR is zero:
1. if backoff == zero then
2.   do broadcast RTR
3.   wait PTRflag is TRUE or waitTerm
4. else wait backoff until zero
5. end if

Received RTR:
1. if countRTR > zero then
2.   do countRTR += 1 reading time
3.   do wait during reading time
4. while waitTime != 0
5.     do wait during reading time
6.   end while
7.   do countRTR -= 1
9. else
```
\begin{algorithm}
\begin{algorithmic}[1]
\State \textbf{do} backoff counter stop
\State \textbf{do} broadcast PTR
\State \textbf{do} \text{countRTR} = 1
\State \textbf{do} \text{waitTime} = reading time
\State \textbf{do} wait during reading time
\State \textbf{end if}
\State \textbf{if} \text{countRTR} == zero \textbf{then}
\State \textbf{goto} countRTR is zero:
\State \textbf{else} \textbf{goto} Received RTR:
\State \textbf{Received PTR}:
\State \textbf{1.} \textbf{waiting backoff}
\State \textbf{2.} \textbf{if} \text{countRTR} == zero \textbf{then}
\State \textbf{3.} \textbf{goto} countRTR is zero
\State \textbf{4.} \textbf{else} \textbf{goto} Received RTR:
\State \textbf{5.} \textbf{end if}
\end{algorithmic}
\end{algorithm}

\textbf{III-II-II ECO\textsuperscript{+}: Power Control in ECO}

If the density of distribution of readers is high, it can lead to the channel capture because of mutual RTR packets. The channel capture effect happens when one node continues to "win" the link. This phenomenon is a cause of starvation problem. In this case, power control is one of methods which can mitigate effectively the capture problem. Power control can also offer better throughput because of spatial reuse. Consequently, we have designed ECO\textsuperscript{+} which executes power control in ECO. In fig 10 (a), R1 suffers from adjacent nodes R2, and R3 which have shorter backoff time than R1. At this time, R1 sends power control message to adjacent nodes with the control channel (2.4GHz), and nodes which receive this message turn down transmission power of both of reader and sensor. Fig 10 (b) shows readable nodes after power control.

Sending power control packet is decided by the number of received RTR of R1. In fig 11, our power control algorithm has power-levels from 1 to $L$, and the power at level $i$ is denoted as $P_i$, where $P_i \in P_A \{P_1, P_2, P_3, \cdots, P_L\}$ and $P_L$ is maximum power. Whenever the reader receives consecutively the RTR packets more than threshold value, it goes to \textbf{DEC:} state and sends \textit{pwrCtrlMsg} to adjacent
nodes, then turns down its current power $P_{\text{cur}}$. On the other hand, power control can deteriorate link connectivity due to decreasing transmission power. To avoid this problem, therefore, the reader should increase transmission power again. When the reader fails to send data because of disconnection with adjacent nodes, $\text{transFail}$ is set TRUE. At this time, the reader goes to $\text{INC}$ state and increases its $P_{\text{cur}}$, and then sends $\text{pwrCtrlMsg}$. $\text{ctrlType}$ variable in the received $\text{pwrCtrlMsg}$ decides to increase or decrease transmission power. If $\text{ctrlType}$ is TRUE, this $\text{pwrCtrlMsg}$ is for decreasing power.

$$T_{\text{WAIT}} = R_t + M_t \quad (1)$$
$$R_t = 2(T_i + T_d) = 2(\text{MAX}(RT_{\text{cal}} + 10T_{\text{pre}} + 11.5T_{\text{pre}})) \quad (2)$$
$$M_t = RT_{\text{cal}} + TR_{\text{cal}} \quad (3)$$

where $R_t$ and $M_t$ are reading time and marginal time respectively, $T_i$ is the time from interrogator transmission to tag response, and $T_d$ is the interrogator response time required if a tag is to demodulate the interrogator signal, measured from the last falling edge of the last bit of the tag response to the first falling edge of the interrogator transmission. $RT_{\text{cal}}$ is interrogator to tag calibration symbol. $TR_{\text{cal}}$ is tag to interrogator calibration symbol. Further details are stated in EPC spec [1]. Table 1 shows the characteristics of the node in integrated networks. Simulation is completed when all readers finish reading.

**IV. Performance Evaluation**

In this section, we will describe our simulation environment and present performances of ECO and ECO' through an analysis of simulation results.

**IV-I Simulation Environment**

Simulation has implemented using Matlab. In a square of 10 by 10 meters, nodes (Ubicons) are randomly uniform-distributed, while their antennae have omni-directional characteristics and transmission range is set to 2 meters. The number of tags is 1000, and the number of readers (Ubicons) is from 25 to 200. In the simulation, we distribute randomly readers and tags according to uniform-distribution such as fig 12. We double the number of nodes each simulation. We use same topology and repeat 10000 times each simulation for fairness. The probability of traffic follows poisson-distribution. The reading time in ECO is assumed to be

Fig. 12. The distribution of reader and tags in the simulation: blue circles are readers (Ubicons), red circles are tags, the green is reading range, and interference range is cyan.

**IV-II Simulation Result**

Fig 13 shows the simulation result of the number of received responses as the number of nodes increases. ECO and ECO' have read more many tags than pure-CSMA. In fig 14, ECO and ECO' take longer total reading completion time than pure-CSMA protocol, as a tradeoff of enhanced stability. Refer to fig 15, the number of failures of ECO and ECO' are lower than pure-CSMA, and it proves that almost every node is able to read tags. The throughput is calculated by

$$\text{System Throughput} = \frac{\sum_i R_i}{\text{Total time}} \quad (4)$$

where, $R_i$ is the number of responses which node $i$ receives from tags. Pure-CSMA has more failures than others. However, total reading complete time of pure-CSMA is also shorter than others, these results also
influence in throughput by equation (4). Although ECO and ECO+ have longer total reading completion time than pure-CSMA, these take better throughput than pure-CSMA in Fig. 16.

Table 1. Characteristics of Ubicon in the simulation environment

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Frequency</td>
<td>2.4GHz (for WSN)</td>
</tr>
<tr>
<td></td>
<td>900MHz (for Tags)</td>
</tr>
<tr>
<td>Communication Range</td>
<td>2m (initial)</td>
</tr>
<tr>
<td>macMinBE</td>
<td>3</td>
</tr>
<tr>
<td>aMaxBE</td>
<td>5</td>
</tr>
<tr>
<td>macMaxCSMABackoff</td>
<td>4</td>
</tr>
<tr>
<td>$T_{\text{WAIT}}$</td>
<td>0.000312</td>
</tr>
<tr>
<td>Scale</td>
<td>100 m²</td>
</tr>
<tr>
<td>Data Size (RTR)</td>
<td>5 byte</td>
</tr>
<tr>
<td>Data Size (PTR)</td>
<td>7 byte</td>
</tr>
<tr>
<td>Control Packet Size (for reader)</td>
<td>11 byte</td>
</tr>
</tbody>
</table>

In every way, ECO+ performance is better than ECO according to density increases, and it proves necessity of power control in dense networks. Moreover, the lowest failure and the best throughput achieved by ECO+ guarantees improvements in the channel capture and spatial reuse.

V. Conclusion

We proposed an enhanced collision avoidance MAC protocol for the integrated RFID-WSN networks. The goal of our proposal is to integrate RFID and WSN, and to avoid reader collisions in this networks. For the network integration, we have designed the integrated network architecture and an integrated module (Ubicon). In the proposed network architecture, sensor-nodes, tags and reader coexist, and the Ubicon arbitrates communication of these components. Ubicon is a sensor-node and a reader interfaced unit. It can communicate not only with sensor-nodes and readers but also can read tags.
The Ubicon enables the integrated network to perform required functions such as ad-hoc, routing, aggregation and environment sensing. These functions enhance capacities of RFID networks. In addition, we implemented ECO for efficient communication without reader collisions. ECO is CSMA-based MAC protocol using two-way handshake, enables several nodes to simultaneously read tags without collision. We also adopted power control mechanism to mitigate the channel capture phenomenon. As a result, when we compared ECO and ECO' with pure-CSMA, throughput was increased by 9.8% and 14.6%, respectively, and the best throughput and the lowest failure shows that ECO' can efficiently break starvation.

References


BIOGRAPHY

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