A Class of Binary Cocyclic Quasi-Jacket Block Matrices

Moon Ho Lee · Subash Shree Pokhrel · Seung Je Choi · Chang Joo Kim

Abstract

In this paper, we present a quasi-Jacket block matrices over binary matrices which all are belong to a class of cocyclic matrices is the same as the Hadamard case and are useful in digital signal processing, CDMA, and coded modulation. Based on circular permutation matrix (CPM) cocyclic quasi block-low-density matrix is introduced in this paper which is useful in coding theory. Additionally, we show that the fast algorithm of quasi-Jacket block matrix.

Key words : Cocyclic Matrices, Jacket Matrices, Hadamard Matrices, Low-Density Matrices.

I. Introduction

The Walsh-Hadamard matrix is widely used for the Walsh representation of the data sequence in image coding and for Hadamard-Walsh orthogonal-sequence generator in code division multiple access (CDMA), spread-spectrum communications,[3][6], coded modulation [7][14]. Their basic functions are sampled Walsh functions which can be expressed in terms of the Hadamard $H_N$ matrices. Using the orthogonality of Hadamard matrices, more general matrices have been developed[9][18]. These matrices are called Jacket matrices. In a general definition, a square matrix $J_{N \times N} = [j_{i,j}]_{N \times N}$ is called Jacket matrix if

$$\frac{1}{N} [j_{i,j}]_{N \times N}$$

for $1 \leq i, j \leq N$, where $C$ is a nonzero constant and $(\cdot)^T$ denotes the transpose of the matrix. Since the inverse of the Jacket matrix can be calculated easily, it is very helpful to employ this kind of matrix in the signal processing[13], encoding[16], multiple-input multiple-output (MIMO) designs[4], sequence design[1], mobile communication[9], and cryptography[18]. In addition, the Jacket matrices are associated with many kind of matrices, such as unitary matrices, and Hermitian matrices which are very important in signal processing, communication (e.g., encoding), orthogonal code design, mathematics and physics.

The outline of the paper is as follows. Section 2 presents a construction of cocyclic quasi-Jacket block matrices over binary matrices and fast algorithm of these quasi-Jacket block matrices. The block matrices may have important application in signal processing, communication, quantum information, etc. A generalized cocyclic quasi-Jacket block low-density matrix based on circular permutation matrix is defined and is constructed in Section 3. Finally, some conclusions are presented in Section 4.

II. Cocyclic Quasi-Jacket Block Matrices

Many interesting matrices, such as Hadamard, WHT, DFT, and its variations matrices are all belong to the class of Jacket family. In many applications, cocyclic matrices are very useful. The definition of cocyclic matrix is in the following.

If $G$ is a finite group of order $v$ and $C$ is a finite Abelian group of order $w$, a (two-dimensional) cocycle is a mapping $\varphi : G \times G \rightarrow C$ satisfying the cocycle equation

$$\varphi(g,h)\varphi(gh,k) = \varphi(g,hk)\varphi(h,k),$$

$\forall g, h, k \in G.$

(1)

This implies

$$\varphi(g,1) = \varphi(1,h) = \varphi(1,1).$$

$\forall g, h \in G$

Then, $\varphi(1,1) = 1$, we call it normalized-cycle[15].

A cocycle associated with the groups $G$ and $C$ is naturally represented as a square matrix of order $v \times v$, whose rows and columns are indexed by the elements of the group $G$ under some fixed ordering and whose entry in position $(g, h)$ is $\varphi(g, h)$. We call such matrices $G$-cocyclic matrices. We represent a $G$-cocyclic matrix as

$$M_\varphi = [\varphi(g, h)]_{g,h \in G}.$$

If the cocycle $\varphi$ is symmetric then $M_\varphi$ is a sy-
metric matrix.

**Definition 2.1:** Let \( [J]_{N \times N} = [(j)_{i,j}]_{N \times N} \) be a block matrix if 
\( [J]_{N \times N} = \left( \frac{1}{C} (j)_{i,j} \right)_{N \times N} \) for \( 1 \leq i, j \leq N \), then the matrix \([J]_{N \times N}\) is a quasi-Jacket block matrix, where \( C \) is a nonzero constant, and \((\cdot)^T\) denotes the transpose of the matrix.

In a quasi-Jacket block matrix, any element of the block Jacket matrices is a common matrix or block matrix which is denoted \((j)_{i,j}\). The elements of the quasi-Jacket block matrix \([J]_{N \times N}\) are substituted with a matrix or even a block matrix. Especially, the inverse of the quasi-Jacket block matrix \([J]_{N \times N}\) may be obtained by the block-wise elements.

Let us consider the following binary matrices:

\[
\alpha \Delta = \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix}, \quad \beta \Delta = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}
\]

(3)

It can be easily checked that

\[
\alpha^2 = \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} \times \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} = I_2,
\]

(4)

\[
\alpha \beta + \beta \alpha = \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} \times \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} + \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} \times \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} = I_4,
\]

(5)

and

\[
\alpha \beta \times \beta \alpha = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \times \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} = I_2,
\]

(6)

i.e.

\[
\alpha = \alpha^{-1}, \quad \beta = \beta^{-1}, \quad \alpha^2 + \beta^2 = O_2,
\]

\[
\alpha^2 = \beta^2 = O_2, \quad \alpha \beta + \beta \alpha = I_2,
\]

\[
(\alpha \beta)^{-1} = \beta \alpha, \quad (\alpha^2 \beta^2)^{-1} = \alpha \beta.
\]

(7)

Furthermore,

\[
\begin{bmatrix} \alpha & \beta \\ \beta & \alpha \end{bmatrix} \times \begin{bmatrix} \beta & \alpha \\ \alpha & \beta \end{bmatrix} = \begin{bmatrix} \alpha \beta + \beta \alpha & \alpha^2 + \beta^2 \\ \beta^2 + \alpha^2 & \beta \alpha + \alpha \beta \end{bmatrix} = \begin{bmatrix} I_2 & O_2 \\ O_2 & I_2 \end{bmatrix}.
\]

(8)

Then we can write,

\[
\begin{bmatrix} \alpha & \beta^{-1} \\ \beta & \alpha \end{bmatrix} = \begin{bmatrix} \beta & \alpha \\ \alpha & \beta \end{bmatrix}.
\]

(9)

Let

\[
[QJ] = \begin{bmatrix} \alpha & \beta \\ \beta & \alpha \end{bmatrix} = \begin{bmatrix} 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \end{bmatrix}.
\]

(10)

Its inverse matrix is

\[
[QJ]^{-1} = [QJ]^T = \begin{bmatrix} \alpha & \beta \\ \beta & \alpha \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 1 \end{bmatrix} = \begin{bmatrix} \alpha & \beta \\ \beta & \alpha \end{bmatrix} = [QJ].
\]

(11)

From definition 2.1, \([QJ]_4\) is a cocyclic quasi-Jacket block matrix. Similarly, we can construct the large quasi-Jacket block matrix based on kronecker product.
Then

\[ [QJ]_4 \times [QJ]_4^{-1} = I_{16}, \]  

(13)

where

\[
\begin{bmatrix}
\beta & \alpha & \beta & \alpha \\
\alpha & \beta & \alpha & \beta \\
\beta & \alpha & \beta & \alpha \\
\alpha & \beta & \alpha & \beta \\
\end{bmatrix}
\begin{bmatrix}
O_1 \\
O_2 \\
O_3 \\
O_4 \\
\end{bmatrix}
= [QJ]^T,
\]

The matrix \([QJ]_4^{-1}\) can be written as follows:

\[
[QJ]_8^{-1} = \begin{bmatrix}
[QJ]_4^{-1} & [QJ]_4^{-1} & [QJ]_4^{-1} & \cdots & [QJ]_4^{-1} \\
\end{bmatrix},
\]

(14)

\[
\begin{bmatrix}
1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 \\
0 & 1 & 1 & 1 & 0 & 1 & 1 & 1 \\
1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 \\
1 & 1 & 0 & 1 & 1 & 1 & 0 & 1 \\
0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\
\end{bmatrix}
\]

(15)

This is also cyclic quasi-Jacket block matrix according to definition 2.1.

Definition 2.2: Let \([QJ]_4\) be defined by (2.5). For \(m \geq 2\), let,

\[
[QJ]_m = [QJ]_4 \otimes [QJ]_4 \otimes \cdots \otimes [QJ]_4,
\]

(17)

\[
= \prod_{r=0}^{m} ([I]_{4^{-1}} \otimes [QJ]_4 \otimes [I]_{4}),
\]

when \(m=2\)

\[
= \prod_{r=0}^{m} ([I]_{4^{-1}} \otimes [QJ]_4 \otimes [QJ]_4).
\]

(16)

The signal flow graph corresponding to (17) is shown in Fig. 1. It requires 16 additions and 32 multiplications for computation.

In general, the computational complexity of the fast quasi-Jacket block fast algorithm is shown in Table 1. In \(GF(2)\) the matrix order \(N=2^m\) requires \(N(\log_2 N - 1)\) additions and \(2N(\log_2 N - 1)\) multiplications.

III. Construction of Cyclic Quasi-Jacket Block Low-Density Matrices Based on Circular Permutation Matrix

In this section, we present low density cyclic quasi block Jacket matrices which are over \(GF(2)\), i.e., binary matrices. These types of low density quasi block Jacket matrices are useful in coding theory such
Table 1. Computational complexity of fast algorithm.

<table>
<thead>
<tr>
<th></th>
<th>Directed Computation</th>
<th>Proposed $N=2^n$</th>
</tr>
</thead>
<tbody>
<tr>
<td>Additions</td>
<td>$N(N-1)$</td>
<td>$N\log_2 N-1$</td>
</tr>
<tr>
<td>Multiplications</td>
<td>$N^2N$</td>
<td>$2N\log_2 N-1$</td>
</tr>
</tbody>
</table>

LDPC encoding, information theory, and orthogonal code design.

Let $p$ be a prime, and let

$$E^h = \left[ e_{i,j} \right]_{E^p},$$

and

$$e_{i,j} =
\begin{cases} 
1 & \text{for } i = (j + h) \\
0 & \text{otherwise}
\end{cases}$$

where $(j + h) = j + h \mod p$ and $0 \leq i, j, h \leq p - 1$.

The matrices $E^h$, for $0 \leq h < p - 1$ refer to circulant permutation matrices (CFM). It can be easily seen that $(I, E^1, \ldots, E^{p-1})$ form an Abelian group with traditional matrix multiplication and $I = E^0$.

**Example:** Let $p = 5$ have

$$I = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}, \quad E = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix},$$

$$E^2 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}, \quad E^3 = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix},$$

$$E^4 = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix}.$$ (20)

Let

$$\Omega = \begin{bmatrix} 1 & O \\ E^h & I \end{bmatrix}, \quad \Omega^2 = \begin{bmatrix} I & E^{-h} \\ O & I \end{bmatrix},$$

$$\Lambda = \begin{bmatrix} I & O \\ E^h & I \end{bmatrix} \times \begin{bmatrix} I & O \\ E^h & I \end{bmatrix} = I_{2^p}, \quad \Lambda^2 = \begin{bmatrix} I & O \\ E^h & I \end{bmatrix} \times \begin{bmatrix} I & E^{-h} \\ O & I \end{bmatrix} = \Omega^2,$$ (21) (22)

and

$$\Lambda = \Lambda^{-1}, \Omega = \Omega^{-1}, \Lambda^2 = \Omega^2 = I_{2^p},$$

$$\Lambda \Omega + \Omega \Lambda = I_{2^p},$$

$$(\Lambda \Omega)^{-1} = \Omega \Lambda, (\Omega \Lambda)^{-1} = \Lambda \Omega.$$ (23) (24) (25)

Furthermore,

$$\begin{bmatrix} \Lambda & \Omega \\ \Omega & \Lambda \end{bmatrix} \times \begin{bmatrix} \Lambda & \Omega \\ \Omega & \Lambda \end{bmatrix} = \begin{bmatrix} \Lambda \Omega + \Omega \Lambda & \Lambda^2 + \Omega^2 \\ \Omega^2 + \Lambda^2 & \Omega \Lambda + \Lambda \Omega \end{bmatrix} = \begin{bmatrix} I & O \\ O & I \end{bmatrix}.$$ (26)

Let,

$$J = \begin{bmatrix} I & O \\ E^h & I \end{bmatrix} \times \begin{bmatrix} I & O \\ E^h & I \end{bmatrix} = J^T.$$ (27) (28) (29) (30) (31)

where

$$J^{-1} = \begin{bmatrix} I & O \\ \Lambda & \Omega \end{bmatrix} \times \begin{bmatrix} I & O \\ E^h & I \end{bmatrix} = \Lambda^{-1} \Omega^{-1}.$$ (32)

From the Definition 2.1, $J$ is a quasi-Jacket which can
be seen that \( J \) is a cocyclic matrix. Now we consider the density of 0's and 1's in \( 4p \times 4p \) binary matrix denoted by

\[
S_{4p} = \frac{N}{N_0 + N_1}
\]

where \( N_0 \) and \( N_1 \) the numbers of 0's and 1 are respectively. The total number of 0's and 1's in \( 4p \times 4p \) binary matrix is \( 16p^2 \). There are 12I, \( E^{-1} \), and \( E^2 \). Each of I, \( E^{-1} \), and \( E^2 \) contains only p 1's. The total number of 1's is 12p. Thus the density of 1's in \( J \) is

\[
S_{4p} = \frac{12p}{16p^2} = \frac{3}{4p}
\]

(33)

It means that this quasi-Jacket matrix is low density matrix. For example, let \( p=5 \) and \( h=2 \), then size of the matrix will be \( 20 \times 20 \), as shown (34)–(36),

\[
[J]_{20} = [J]_4 \otimes [I]_4,
\]

\[
= ([I]_4 \otimes [J]_4)([J]_4 \otimes [I]_4)
\]

(34)

Remark: \( J \) is a (3,3)-low density matrix, i.e., each row and each column of \( J \) contains exactly three 1's. The quasi-Jacket block matrix of order \( 20 \times 20 \) contains four 1's in the corner of the matrix, we can say it as a cycle-4 or girth-4. Generally, short cycles have negative impact on the decoding algorithm for LDPC \[19\]. To use these kinds of matrices in LDPC, once should remove the cycle-4 or girth-4 of the matrices.

IV. Conclusion

In this paper, we have presented cocyclic quasi-Jacket block matrices over binary matrices which all are belong to a class of cocyclic matrices. We have

\[
\begin{bmatrix}
1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0
\end{bmatrix}
\]

(35)

\[
J^{-1}
\]

(36)

also shown how to construct low-density cocyclic quasi-Jacket matrices based on circular permutation matrix. These cocyclic quasi-Jacket block matrices and low-density matrices are useful in digital signal processing, CDMA, coding theory, orthogonal code design, and coded modulation \[1\]–\[3\]\[16\].

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