Encryption Algorithm using Polyline Simplification for GIS Vector Map

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ABSTRACT

Recently, vector map has developed, used in many domains, and in most cases vector map data contains confidential information which must be kept away from unauthorized users. Moreover, the manufacturing process of a vector map is complex and the maintenance of a digital map requires substantial monetary and human resources. This paper presents the selective encryption scheme based on polyline simplification methods for GIS vector map data protection to store, transmit or distribute to authorized users. Main advantages of our algorithm are random vertices and transformation processes but it still meets requirements of security by random processes, and this algorithm can be implemented to many types of vector map formats.

Key words: GIS Vector Map, Selective Encryption, DWT, Randomization.

1. INTRODUCTION

Vector map is created and developed by the merging system of cartography, statistical analysis, and database technology based on vector model [1, 2]. Vector map stores and manages all kinds of the geographic information data as geometric factor, topology and metadata by vector data.

Vector data provide a way to represent real world features within the GIS environment because vector data has advantages as need a small space or place for storage data; easily makes connection between topology and network; has a high spatial resolution and graphic representation spatial data closely likes handed map; easily for making projection and coordinates transformation. But the producing process is considerably complex and the maintenance of a digital map requires substantial monetary. So vector map is necessary to be protected and prevent illegal duplication and distribution of it. Moreover, applications of vector map in military domain require the high security, and must be kept away from unauthorized users. So vector map is necessary to be protected and prevent illegal duplication and distribution of it. In proposed algorithm, vector map is separated to select polyline/polygon layer. The three simplification algorithms are used to define the feature points in each object. After that, the randomly vertices of objects are selected based on the ratio of the feature points. Finally, it is randomized by randomization values and encrypted by using DWT transform.

Our paper is organized as follows. In section 2,
we discuss the related works and in section 3, we explain the proposed selective encryption algorithm in detail. Then, in section 4 we perform experiments and discuss about the experimental results, evaluate the performance of algorithm. Finally, we conclude this paper in section 5.

2. RELATED WORKS

According to the recent growth of network digital media, data are needed to protect from distribution and illegal copying. Many approaches are researched for this issues; these include authentication, encryption, and time stamping.

Digital watermarking has been researched to solve the issues in vector map since 2000s. The traditional method embed secret information in some locations of vector map; [3] proposes certain rules, they help to select a set of coordinates of vertices, and editing them by using a certain range of precision; [4] modify coefficients in the frequency domain to complete watermarking hiding that is one of the important solution of watermarking algorithm, but the disturb to the vector map content is also existing, the resistance performance to data fitting, interpolation, scaling is poor. Due to the vector map high precision requirement, traditional watermarking algorithm can't meet the demand of practical application.

The full encryption algorithms usually encrypt all components of original data to change whole data. Wu et al. [5] consider characteristics of the storage, parameters and initial values of chaotic map. After that, he proposed a new compound encryption method, this process is not available to any type formats of data and object indexing. Li et al. [6] selected the vector dataset in external Oracle DBMS for encryption, and he used standard cryptography DES combining with an R-Tree spatial index. This algorithm encrypts the spatial index when the GIS dataset is transmitted to the client and designs the key management of public and private keys on a PKI system. In this process, the key length is short so it can not keep data on the DBMS with high security. Dakrouy et al. [7] also described better the encrypting algorithm, AES and RSA cryptography are combined along with watermarking method that used in internet online service. This algorithm encrypts all parts of a shape-file using 256 bit for private key of a block cipher AES. But, this algorithm uses whole shape-files for encryption and they do not consider important features, it is taking a long time to handle.

The general approach is to separate the content into two parts. The first part is the public part, it is left unencrypted and made accessible to all users. The second part is the protected part; it is encrypted. Only authorized users have access to protected part. One important feature in selective encryption is to make the protected part as small as possible. The control factor (threshold value) helps to control encrypted data ratio in algorithm, by this way we can make security better. In the main, perceptual process is created based on partial encryption algorithms.

For vector map perceptual encryption, B.-J Jang et al. [8] presented a method, which encrypts parameters mean point and direction of mini coding objects in compression domain by XOR operator. The aim of this method is to select parameters after compression to encrypt by XOR operator. Giao et al. [9] select all vertices in a complex layer and they randomize them before to encrypt in DCT domain, they did not consider the important part in each layer. For these reasons, a new algorithm is proposed in selective encryption topic for multimedia applications, transmission and storage of vector

![Diagram](image_url)

**Fig. 1.** The concept of selective encryption,
map data.

3. PROPOSED ALGORITHM

In our paper, the proposed method selects randomly vertices for encryption based on simplification algorithms. In Fig. 2, we show the schematic of algorithm, and the step-by-step is explained in detail:

The vector map data includes layers and each layer includes geometry feature which is illustrated by geometric objects, and attribute features. Attribute data describes specific map features but is not inherently graphic. The geographic information used a set of coordinates to provide position and shapes of objects as point, polyline and polygon. Thus, we consider the content of layers in a map consists of two parts. The first part includes insignificant information as text and annotation. The second part is the geometric data needed to protect, includes geometric objects as point, polyline and polygon.

Moreover, the point only uses a pair of coordinates to represent simple, small and zero-dimensional objects in the real on the map while polylines and polygons use a set of coordinate to represent complex structure and huge objects. Therefore, polylines and polygons are targets to select for vector map security.

3.1 Backbone of object

As you know, a layer is set of objects and an object is a set of vertices \( O_j = \{v_{ij}|k \subseteq |O_j|\} \) with \( |O_j| \) is the cardinality of the object. The object is separated to \( m \) groups \( G_k = \{G_k|k \subseteq |m|\}; m = |O_j|/n; n \) is the total of vertices in one group. Object’s backbone is set of vertices with a vertex is average point between two continuous vertices in a group, as shown in Fig. 3. Therefore, backbone \( B = \{B_k|k \subseteq |m|\} \) with \( B_k \) is a backbone section in group \( G_k \).

3.2 Breakpoint definition

We use the scale factor \( \alpha \) for controlling the simplification quantity and for generating breakpoints with different ratios. Thus, we set parameters for the three simplification algorithms using the scale factor and initial parameters:

\[
\delta = \delta_0(1+\alpha) \quad \text{for DP} \quad (1) \\
\epsilon = \epsilon_0(1+\alpha) \quad \text{for SF} \quad (2) \\
\gamma = \gamma_0(1+\alpha) \quad \text{for LA} \quad (3)
\]

where \( \delta_0, \epsilon_0, \) and \( \gamma_0 \) are the initial distance thresholds for DP, SF, and LA. The scale factor \( \alpha \) controls the number of breakpoints. If \( \alpha = -1 \), a backbone is not simplified. If \( \alpha \) increases from \(-1\), the total number of breakpoints also increase.

Given a backbone \( B \) of object, let us consider

![Fig. 2. The proposed encryption process of vector map data.](image)
that $\mathcal{B}$ consists of $N$ vertices, $B = \{v_{b1}, v_{b2}, \ldots, v_{bN}\}$, and a vertex has two coordinate values $v_{bi} = (x_{bi}, y_{bi})$. We generate three simplified backbones of backbone $\mathcal{B}$ using the DP, SF, and LA algorithms: $B_{DP}, B_{SF}$, and $B_{LA}$.

Breakpoints $\mathcal{B}'$ for a backbone $\mathcal{B}$ as vertices that exist in three simplified backbone:

$$B = B_{DP} \cap B_{SF} \cap B_{LA} = \{v_{bi} | k \in [1, N_B] \}$$  \hspace{1cm} (4)

Thus, breakpoints are common vertices in three simplified backbone of a backbone. This mean that they represent the main shape of the polyline and should be kept simplified. $N_B$ is the number of vertices in the breakpoints. In the next step, breakpoints is used to select vertices group for encryption.

### 3.3 Vertices encryption

In mathematics, a chaotic map is a map that exhibits some sort of chaotic behavior, by following:

$$x_{k+1} = nx_k(1-x_k)$$  \hspace{1cm} (5)

We use the user key to create random coefficients $R_i$ and key values for a layer. These values are created randomly the first values by SHA-512 process using key given by user [9], 512 bit key size. And other values are generated by using Chaotic map as equation (5).

In group $G_k$, if the total number of vertices of backbone section $R_i$ bigger than $N_B / m$ ( $N_B$ is the number of breakpoints, $m$ is the number of groups), we select all vertices in the selected group $G_k$ and encrypt them by steps, as follow:

- We arrange all the first vertex of the selected groups in a layer into 1D-array $V = \{v_{i1}, v_{i2}, \ldots, v_{in}\}$.
- These vertices are randomized by multiplying them using random coefficients $R_i$.
- After that, we apply DWT-$k$ level to get DWT coefficients. In DWT domain, we encrypt level $k$ coefficient and inverse DWT-$k$ level. We get $V = \{v'_{i1}, v'_{i2}, \ldots, v'_{in}\}$.
- With the selected group $G_k$, other vertices are changed following the variation of the first vertex in this group by equation:

$$v'_{n^*(k-1)+i} = v_{n^*(k-1)+i} + (-1)^{n^*(k-1)+i} v'_{ik1}$$  \hspace{1cm} (6)
3.4 Decryption process

The reverse process is applied to decrypt the encrypted map. To perform decryption, after we calculate thresholds in a layer, we use the decryption block to decrypt the encrypted significant objects. If correct key is employed at the time of decryption, then the decrypted map would by a replica of the original map. In the case of an incorrect decryption key, the output map is vary.

4. EXPERIMENTAL RESULTS

4.1 Vertices selection

Our method select randomly vertices based on different precisions by adjusting the scale factor $\alpha$ in Eq. (1)–(3) (that determines the quantity of simplification). $\alpha = -1$ means that our method by passes the simplification. $\alpha = 0$ means that our method simplifies a map using the initial parameters $(\delta_0, \epsilon_0, \gamma_0)$ of the DP, SP, and LA algorithms. Fig. 4 shows the ratio of breakpoint number and encrypted vertices calculated by formula: number of breakpoints/ number of vertices of object’s backbone.

4.2 Visualization

Experimental results Fig. 5 and Fig. 6 show the proposed algorithm changes whole maps. The proposed method has much lower computational complexity than AES or DES because we only select and encrypt some vertices in objects.

4.3 Algorithm comparison

The proposed method is compared with two existing algorithms Giao [10] and Bang [11]. In these.

Fig. 4, Ratio of breakpoints and encrypted vertices.
papers, first author encrypted all data in original file and he did not consider important part in each layer. The second author selected all vertices of the significant objects (defined based on threshold value) in a layer for encryption, it seem quite easily to attack, because he defined threshold value very simple lead to the algorithm is so weak. Our method select randomly some vertices (that is defined by owner based on scale factor a and three initial parameters) for encryption, attacker will be very difficult to determine the encrypted data and predict key. Fig. 7 shows the ratio of encryption data in comparison with two existing methods.

5. CONCLUSION

In our paper, we created a new method which aim to reduce ratio of encrypted data in GIS vector map but still assure performance and high security. This considers how to select randomly vertices of object in a layer and encrypts randomly components by simplification algorithm and Chaotic-map in DWT domain. This process encrypt some level-k DWT coefficients lead to change whole map. We also confirm that Human perception do not see any information in encrypted map, poor error in decryption step, computation time is less than it in the existing methods, high security and a large amount of GIS vector map data can be protected by this algorithm.

REFERENCE


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